Quantum Query Algorithms for Automorphisms of Galois Groups

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Abstract—In this paper we study quantum query complexity of exactly (with probability 1) deciding the parity of n-permutations of numbers (from 0 to n-1). We show that for this non-Boolean problem use of quantum complexity techniques gives quite strong results as it does for other, but Boolean problems. We use Galois Theory to find new problems where quantum query algorithms are more efficient than deterministic ones.

Keywords—quantum computing; quantum query algorithms; galois groups; permutations.

1. QUANTUM QUERY ALGORITHMS

Query algorithms are frequently used to compute Boolean functions. As an input for the query algorithm a black box oracle is used. It contains the values of the variables $x_1, x_2, \ldots, x_n$ for known Boolean function $f(x_1, \ldots, x_n)$

The complexity of the query algorithm is equal (by definition) to the number of the queries asked to the black box oracle. In deterministic query algorithms the next query is chosen uniquely depending only on answers previously received from the black box oracle. Probabilistic query algorithms allow randomization of the process of computation i.e. the choice of next query can depend on some randomness element.

Quantum query algorithms are formally defined in [1]. For the sake of this paper we will introduce some background and parts that are necessary for our research.

Quantum query algorithm has a finite number of states. In each state it is allowed to make a query to the black box oracle and to change states in accordance to used algorithm. There are two types of algorithm steps: query operations and unitary transformations, these steps are performed in alternating order – one after the other. In the query operations step the states are divided into subsets corresponding to each query then an allowed quantum-operations step the states are divided into subsets in alternating order – one after the other. In the query

In quantum query algorithms with a promise is notion of promise for quantum algorithm. The domain of correctness in quantum query algorithms with a promise is explicitly restricted. We are only interested in values of the algorithm that follow this restriction and do not look at other values. For instance, in this paper there is a promise on all the query algorithms that the calculated function represents a permutation.

There are many recent papers studying query algorithms that are computing Boolean functions. For reference see the survey by Buhrman and de Wolf [1].

In this paper we consider a class of functions $f(x_1, \ldots, x_n)$ that are more general.

We consider a particularly interesting case of the domain $\{0,1,2, \ldots, n\}$-permutations.

For instance, $x_1 = 4, x_2 = 3, x_3 = 2, x_4 = 0, x_5 = 1$

represents a 5 symbol $\{0,1,2,3,4\}$permutation. It can be described as permutation 43201.

For this special case of domain the functions $f: \{0,1,2, \ldots, n\}^n \rightarrow \{0,1\}$ can be seen as representing properties of permutations. For instance, the function

represents the property of 3 element permutation to be even (The other opposite property of permutation would be to be odd). Deciding the parity of permutation is defined as finding whether this permutation is even or odd.

Quantum query algorithms where the black box contains a permutation were first considered by K.Iwama and
Galois Theory is described in many textbooks [2]. Since then several authors have compared quantum and deterministic query algorithms for permutations but they were not able to find examples of properties of permutations where the complexity of quantum query algorithm is considerably smaller than the complexity of every deterministic query algorithm for the same problem.

II. GALOIS GROUPS OF POLYNOMIALS

Roots of a given polynomial can be connected by different algebraic equations. For example, for two roots, \( A \) and \( B \), the equation \( A^2 + 5B^3 = 7 \) can hold. Galois Theory considers roots with the property that for all the algebraic properties satisfied by the roots they still hold after permutations (or rearrangements) of the roots.

It should be noted that only algebraic equations with rational coefficients are considered. (We could specify a certain field that should contain the coefficients but, for the sake of examples below, we will use the field of rational numbers.) All permutations for a given equation satisfying this property form a permutation group, called the Galois group of the polynomial (over the rational numbers). Consider the following examples that will illustrate this point:

The roots of quadratic equation \( x^2 - 4x + 1 = 0 \) can be found by using the quadratic formula. The two roots are \( A = 2 + \sqrt{3} \) and \( B = 2 - \sqrt{3} \). Some of the algebraic equations that are satisfied by \( A \) and \( B \) include \( A + B = 4 \) and \( AB = 1 \).

Obviously, if we exchange \( A \) and \( B \), in either of these equations, we obtain another true statement. For example, the equation \( A + B = 4 \) becomes \( B + A = 4 \).

To fact that this holds for every possible algebraic equation with rational coefficients relating the \( A \) and \( B \) values above (in any such equation, swapping \( A \) and \( B \) yields another true equation) is less obvious but true. The proof is done by using the theory of symmetric polynomials.

There are algebraic equations that relate \( A \) and \( B \) that are false when \( A \) and \( B \) are exchanged (for example equation \( A - B = 2\sqrt{3} = 0 \)), however, this equation is not considered as it has the coefficient \(-2\sqrt{3}\) which is irrational.

We can see that the Galois group of the polynomial \( x^2 - 4x + 1 \) consists of two permutations: the transposition permutation which exchanges \( A \) and \( B \) and the identity permutation which leaves \( A \) and \( B \) untouched. It forms a cyclic group of order two.

Galois Theory is described in many textbooks [2].

III. RESULTS

If a polynomial can be factored into monomials it is said that polynomial splits over a field \( K \). For example, \( x^2 + 3x + 2 \) splits in the rationals, with factors \((x+2)\) and \((x+1)\). Let \( S \) be a set of polynomials taken from \( K[x] \).

The splitting field for \( S \) is the smallest extension that splits all the polynomials in \( S \). Call this extension \( F/K \). If \( F/K \) is a field extension, its Galois group is the group of automorphisms of \( F \) that fix \( K \). If \( p(x) \) is an irreducible polynomial with coefficients in \( K \), it also has a Galois group, namely the Galois group of its splitting field. The splitting field \( F/K \) is the smallest field \( F \) that includes all the roots of \( p(x) \).

Hence if \( p(x) \) is a polynomial with rational coefficients, then there is a smallest subfield \( L(p) \) of the complex numbers \( \mathbb{C} \) containing both the rationals \( \mathbb{Q} \) and all the roots of \( p(x) \). This is the splitting field of \( p \) over \( \mathbb{Q} \). For the polynomial \( p(x) = (x^2 - 2)(x^2 - 3)(x^2 - 5) \), it turns out that \( L(p) = \mathbb{Q}(\sqrt{2}, \sqrt{3}, \sqrt{5}) \). Now this field contains, besides itself and \( \mathbb{Q} \), fourteen other subfields which form \( \{7, 3, 1\} \) design. The varieties of the design are the seven quadratic sub fields \( \mathbb{Q}(\sqrt{d}) \) where \( d \in \{2,3,5,6,10,15,30\} \), and the blocks are the seven biquadratic subfields \( \mathbb{Q}(\sqrt{d_1}, \sqrt{d_2}) \) where \( d_1, d_2 \) is not a perfect square.

We consider the query complexity of the problem \( P \) for the problem \( P \) accepting permutations in \( P \) with probability \( 1 \), and rejecting permutations not in \( P \) with probability strictly exceeding \( \frac{1}{2} \).

(1) Every deterministic query algorithm for the problem \( P \) needs at least 6 queries.

(2) There exists a quantum query algorithm with 2 queries for the problem \( P \) accepting permutations in \( P \) with probability \( 1 \), and rejecting permutations not in \( P \) with probability strictly exceeding \( \frac{1}{2} \).

(3) There exists a quantum query algorithm with 3 queries for the problem \( P \) accepting permutations in \( P \) with probability \( 1 \), and rejecting permutations not in \( P \) with probability \( 1 \).

ACKNOWLEDGMENT

The research was supported by Project 271/2012 from the Latvian Council of Science.

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