

Regular fuzzy equivalences on multi-mode multi-relational fuzzy networks

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Abstract

In this paper we introduce the concepts of a multi-mode multi-relational fuzzy network and a regular fuzzy equivalence on such a network, and provide procedures for computing the greatest regular fuzzy and crisp equivalences contained in a given tuple of fuzzy equivalences.

Keywords: Multi-mode fuzzy network, regular fuzzy equivalence, fuzzy relation equation, residual of fuzzy relations.

1. Introduction

Social network analysis has originated as a branch of sociology and mathematics which provides formal models and methods for the systematic study of social structures. Social networks share many common properties with other types of networks, and methods of social network analysis are nowadays applied to the analysis of networks in general, including many kinds of networks that arise in computer science, physics, biology, etc., such as the hyperlink structure on the Web, the electric grid, computer networks, information networks or various large-scale networks appearing in nature.

In large and complex networks it is impossible to understand the relationship between each pair of individuals, but to a certain extent, it may be possible to understand the system, by classifying individuals and describing relationships on the class level. For instance, individuals in the same class can be considered to occupy the same position, or play the same role in the network. The main aim of the positional analysis of networks is to find similarities between individuals which have to reflect their position in a network. These similarities have been formalized first by Lorrain and White [17] by the concept of a structural equivalence. Informally speaking, two individuals are considered to be structurally equivalent if they have identical neighborhoods. However, in many situations this concept has shown oneself to be too strong. Weakening it sufficiently to make it more appropriate for modeling social positions, White and Reitz [21] have introduced the concept of a regular equivalence, where two individuals are considered to be regularly equivalent if they are equally related to equivalent others. Regular equivalences also play an important role in the blockmodeling, a method of data reduction which

reduces redundant elements of a network to yield a simplified model of relationships between types of elements (cf. [1, 7, 9]).

Regular equivalences have been studied mainly in the context of one-mode social networks, consisting of a single set of entities and ties between these entities of the same or different types. To a lesser extent, regular equivalences and blockmodeling have been studied in the context of two-mode networks, which consist of two sets of entities and ties from one to another set (cf. [1, 4, 7, 8, 9, 18, 19, 20]). However, in real situations we often encounter much more complex networks consisting of multiple sets of entities and ties inside and between some of them. The main aim of this paper is to provide a general mathematical model for the study of such complex networks, which are called multi-mode networks, to introduce the concept of a regular equivalence on these networks, and to provide a procedure for computing these regular equivalences. It should be noted that a similar kind of networks, called multilevel networks, has been recently studied in [22]. In addition, we consider an even more general case, also very common in real situations, where the ties between entities are fuzzy. Such fuzzy social networks have been studied in [6, 10, 11, 12, 14, 15, 16].

The methodology used here for computing regular fuzzy equivalences on multi-mode fuzzy networks is based on methods for finding the greatest solutions of so-called weakly linear systems of fuzzy relation inequalities and equations, developed in [14, 15, 16], where residual of fuzzy relations play a key role. It is worth noting that this methodology has been previously shown to be very efficient in solving some fundamental problems of the theory of fuzzy automata, such as the reduction of the number of states and the problems of equivalence, simulation and bisimulation.

After this introductory section, in Section 2 we define the basic concepts of the theory of fuzzy sets and fuzzy relations that are used in further work. Then in Section 3 we define a multi-mode multi-relational fuzzy network and the concept of a regular fuzzy equivalence on such a network, and present the main results of the paper which provide a procedure for computing the greatest regular fuzzy and crisp equivalences on such networks.

The results obtained in this paper generalize the related results concerning both one-mode and two-mode fuzzy networks.

2. Preliminaries

In the paper we use complete residuated lattices as structures of membership values. A *residuated lattice* is an algebra $\mathcal{L} = (L, \wedge, \vee, \otimes, \rightarrow, 0, 1)$ such that

- (L1) $(L, \wedge, \vee, 0, 1)$ is a lattice with the least element 0 and the greatest element 1,
- (L2) $(L, \otimes, 1)$ is a commutative monoid with the unit 1,
- (L3) \otimes and \rightarrow form an *adjoint pair*, i.e., they satisfy the *adjunction property*: for all $x, y, z \in L$,

$$x \otimes y \leq z \Leftrightarrow x \leq y \rightarrow z. \quad (1)$$

If, additionally, $(L, \wedge, \vee, 0, 1)$ is a complete lattice, then \mathcal{L} is called a *complete residuated lattice*.

The operations \otimes (called *multiplication*) and \rightarrow (called *residuum*) are intended for modeling the conjunction and implication of the corresponding logical calculus, and supremum (\vee) and infimum (\wedge) are intended for modeling the existential and general quantifier, respectively. An operation \leftrightarrow given by

$$x \leftrightarrow y = (x \rightarrow y) \wedge (y \rightarrow x), \quad (2)$$

called *biresiduum* (or *biimplication*), is used for modeling the equivalence of truth values. It can be easily shown that with respect to \leq , \otimes is isotonic in both arguments, \rightarrow is isotonic in the second and antitonic in the first argument. For other properties of complete residuated lattices we refer to [2, 3].

The most studied and applied structures of membership values, defined on the real unit interval $[0, 1]$ with $x \wedge y = \min(x, y)$ and $x \vee y = \max(x, y)$, are the *Lukasiewicz structure* ($x \otimes y = \max(x + y - 1, 0)$) and $x \rightarrow y = \min(1 - x + y, 1)$, the *Goguen (product) structure* ($x \otimes y = x \cdot y$, $x \rightarrow y = 1$ if $x \leq y$ and $= y/x$ otherwise) and the *Gödel structure* ($x \otimes y = \min(x, y)$, $x \rightarrow y = 1$ if $x \leq y$ and $= y$ otherwise). Another important set of truth values is the set $\{a_0, a_1, \dots, a_n\}$, $0 = a_0 < \dots < a_n = 1$, with $a_k \otimes a_l = a_{\max(k+l-n, 0)}$, $a_k \rightarrow a_l = a_{\min(n-k+l, n)}$. A special case of the latter algebras is the two-element Boolean algebra of classical logic with the support $\{0, 1\}$. The only adjoint pair on this Boolean algebra consists of the classical conjunction and implication operations. This structure of truth values is called the *Boolean structure*.

In the sequel \mathcal{L} will be a complete residuated lattice. A *fuzzy subset* of a set A over \mathcal{L} , or simply a *fuzzy subset* of A , is any mapping from A into L . Ordinary crisp subsets of A are considered as fuzzy subsets of A taking membership values in the set $\{0, 1\} \subseteq L$. Let f and g be two fuzzy subsets of A . The *equality* of f and g is defined as the usual equality of mappings, i.e., $f = g$ if and only if $f(x) = g(x)$, for every $x \in A$. The *inclusion* $f \leq g$ is also defined pointwise: $f \leq g$ if and only if $f(x) \leq g(x)$, for every $x \in A$. Endowed with this partial order the set L^A of all fuzzy subsets of A forms a complete lattice, in which the meet (intersection) $\bigwedge_{i \in I} f_i$ and the join (union) $\bigvee_{i \in I} f_i$ of an

arbitrary family $\{f_i\}_{i \in I}$ of fuzzy subsets of A are mappings from A into L defined by

$$\left(\bigwedge_{i \in I} f_i \right) (x) = \bigwedge_{i \in I} f_i(x), \quad \left(\bigvee_{i \in I} f_i \right) (x) = \bigvee_{i \in I} f_i(x),$$

for all $x \in A$.

A *fuzzy relation* between sets A and B (in this order) is any fuzzy subset of $A \times B$, and the equality, inclusion (ordering), joins and meets of fuzzy relations are defined as for fuzzy sets. The set of all fuzzy relations between A and B will be denoted by $L^{A \times B}$. In particular, a fuzzy relation on a set A is any fuzzy subset of $A \times A$, and the set of all fuzzy relations on A is denoted by $L^{A \times A}$. The *reverse* or *inverse* of a fuzzy relation $\alpha \in L^{A \times B}$ is a fuzzy relation $\alpha^{-1} \in L^{B \times A}$ defined by $\alpha^{-1}(b, a) = \alpha(a, b)$, for all $a \in A$ and $b \in B$. A crisp relation is a fuzzy relation taking values only in the set $\{0, 1\}$, and if α is a crisp relation between A and B , then the expressions " $\alpha(a, b) = 1$ " and " $(a, b) \in \alpha$ " will have the same meaning. For a fuzzy relation $\alpha \in L^{A \times B}$, the crisp relation α^c given by

$$\alpha^c(a, b) = \begin{cases} 1 & \text{if } \alpha(a, b) = 1, \\ 0 & \text{otherwise} \end{cases},$$

is called the *crisp part* of α (it is also known as the *1-cut* of α). According to the above mentioned convention, we also write

$$\alpha^c = \{(a, b) \in A \times B \mid \alpha(a, c) = 1\}.$$

For non-empty sets A , B and C , and fuzzy relations $\alpha \in L^{A \times B}$ and $\beta \in L^{B \times C}$, their *composition* is a fuzzy relation $\alpha \circ \beta \in L^{A \times C}$ defined by

$$(\alpha \circ \beta)(a, c) = \bigvee_{b \in B} \alpha(a, b) \otimes \beta(b, c), \quad (3)$$

for all $a \in A$ and $c \in C$. When the underlying sets are finite, fuzzy relations can be interpreted as matrices with entries in L and the composition of fuzzy relations can be interpreted as a kind of matrix product. It is easy to verify that the composition of fuzzy relations is associative and distributive over unions (joins) of fuzzy relations.

Let A , B and C be non-empty sets, $\lambda \in L^{A \times B}$, $\mu \in L^{B \times C}$ and $\eta \in L^{A \times C}$. The *right residual* of η by λ is a fuzzy relation $\lambda \backslash \eta \in L^{B \times C}$ defined by

$$(\lambda \backslash \eta)(b, c) = \bigwedge_{a \in A} \lambda(a, b) \rightarrow \eta(a, c), \quad (4)$$

for all $(b, c) \in B \times C$, and the *left residual* of η by μ is a fuzzy relation $\eta / \mu \in L^{A \times B}$ defined by

$$(\eta / \mu)(a, b) = \bigwedge_{c \in C} \mu(b, c) \rightarrow \eta(a, c), \quad (5)$$

for all $(a, b) \in A \times B$. It is not hard to verify that the following *residuation property* (in some sources

called the *adjunction property*) holds for arbitrary $\lambda \in L^{A \times B}$, $\mu \in L^{B \times C}$ and $\eta \in L^{A \times C}$:

$$\lambda \circ \mu \leq \eta \Leftrightarrow \mu \leq \lambda \backslash \eta \Leftrightarrow \lambda \leq \eta / \mu. \quad (6)$$

A fuzzy relation $\alpha \in L^{A \times A}$ is *reflexive* if $\Delta_A \leq \alpha$ (where Δ_A is the crisp equality on A), *symmetric* if $\alpha^{-1} \leq \alpha$, and *transitive* if $\alpha \circ \alpha \leq \alpha$. A reflexive, symmetric and transitive fuzzy relation is called a *fuzzy equivalence*, whereas a reflexive and transitive fuzzy relation is called a *fuzzy quasi-order*.

Hereinafter, \mathbb{N} denotes the set of natural numbers (without zero), $\mathbb{N}^0 = \mathbb{N} \cup \{0\}$, and $[1, n]$ denotes the set of first n natural numbers, for each $n \in \mathbb{N}$. For a given family of fuzzy relations \mathcal{R} , by $\mathcal{L}(\mathcal{R})$ we denote the subalgebra of \mathcal{L} generated by all membership values taken by fuzzy relations from \mathcal{R} .

Let $n \in \mathbb{N}$ and let A_1, \dots, A_n be a collection of non-empty sets. The set $L^{A_1 \times A_1} \times \dots \times L^{A_n \times A_n}$ of all n -tuples of fuzzy relations on A_1, \dots, A_n , respectively, is ordered pointwise, in the following way: $(\alpha_1, \dots, \alpha_n) \leq (\beta_1, \dots, \beta_n)$ if and only if $\alpha_l \leq \beta_l$, for each $l \in [1, n]$. Following the terminology that is used for fuzzy sets and relations, we will say that $(\alpha_1, \dots, \alpha_n)$ is contained in $(\beta_1, \dots, \beta_n)$.

A partially ordered set P is said to satisfy the *descending chain condition*, shortly *DCC*, if each descending sequence of elements of P stabilizes, i.e., if $\{a_k\}_{k \in \mathbb{N}}$ is a sequence of elements of P such that $a_{k+1} \leq a_k$, for each $k \in \mathbb{N}$, then there exists $n \in \mathbb{N}$ such that $a_n = a_{n+m}$, for each $m \in \mathbb{N}$.

For more information on fuzzy sets and fuzzy relations we refer to [2, 3].

3. Regular fuzzy equivalences

In this section we first define a multi-mode multi-relational fuzzy network and the concept of a regular fuzzy equivalence on such a network.

Given $n \in \mathbb{N}$, a collection A_1, \dots, A_n of non-empty sets and a non-empty set $J \subseteq [1, n] \times [1, n]$, and let $\{I_{j,k}\}_{(j,k) \in J}$ be also a collection of non-empty sets. We require that for each $j \in [1, n]$ there is $k \in [1, n]$ such that $(j, k) \in J$ or $(k, j) \in J$. For any $(j, k) \in J$ let $\{R_i^{j,k}\}_{i \in I_{j,k}}$ be a family of non-empty fuzzy relations between A_j and A_k , and set

$$\mathcal{R} = \{R_i^{j,k} \mid (j, k) \in J, i \in I_{j,k}\}.$$

The system $\mathcal{N} = (A_1, \dots, A_n, \mathcal{R})$ is called a *multi-mode multi-relational fuzzy network*, or more specifically, an *n-mode multi-relational fuzzy network*. We usually omit the adjective “multi-relational”. The sets A_1, \dots, A_n are called the *components* or *modes* of the network \mathcal{N} . Note that the above families of fuzzy relations are specified for some pairs of components, but not necessarily for all pairs. It is evident that J is the set of all pairs $(j, k) \in [1, n] \times [1, n]$ for which a non-empty family of non-empty fuzzy relations between A_j and A_k is specified, and this family is indexed by the set $I_{j,k}$. If for some $j \in [1, n]$ there

is no any $k \in [1, n]$ such that a non-empty family of non-empty fuzzy relations between A_j and A_k or between A_k and A_j is specified, then A_j is isolated and there is no sense to consider it. For this reason we have introduced the above requirement for the set J .

Observe that the multi-mode network \mathcal{N} is a complex system which consists of one-mode fuzzy networks $\mathcal{N}_j = (A_j, \{R_i^{j,j}\}_{i \in I_{j,j}})$, for those $j \in [1, n]$ for which $(j, j) \in J$, and two-mode fuzzy networks $\mathcal{N}_{j,k} = (A_j, A_k, \{R_i^{j,k}\}_{i \in I_{j,k}})$, for those $(j, k) \in J$ for which $j \neq k$. Therefore, for $n = 1$ this definition gives a one-mode fuzzy network, and for $n = 2$ and $J = \{(1, 2)\}$ it gives a two-mode fuzzy network.

For any $l \in [1, n]$ let α_l be a fuzzy equivalence on A_l such that for all $(j, k) \in J$, $i \in I_{j,k}$ the following is true:

$$\alpha_j \circ R_i^{j,k} = R_i^{j,k} \circ \alpha_k. \quad (7)$$

Then the n -tuple $(\alpha_1, \dots, \alpha_n)$ is said to be a *regular fuzzy equivalence* on the network \mathcal{N} . In addition, if $\alpha_1, \dots, \alpha_n$ are crisp equivalences, then we say that $(\alpha_1, \dots, \alpha_n)$ is a *regular crisp equivalence* on \mathcal{N} . It is clear that $(\alpha_1, \dots, \alpha_n)$ is a regular fuzzy equivalence on \mathcal{N} if and only if α_j is a regular fuzzy equivalence on the one-mode network \mathcal{N}_j , for each $j \in [1, n]$ such that $(j, j) \in J$, and the pair (α_j, α_k) is a regular fuzzy equivalence on the two-mode network $\mathcal{N}_{j,k}$ for each $(j, k) \in J$ for which $j \neq k$. The role of the regular fuzzy equivalence $(\alpha_1, \dots, \alpha_n)$ can be understood as follows: for any $j \in [1, n]$ the fuzzy equivalence α_j identifies and classifies similar entities in the component A_j , entities that occupy the same position or play the same role in the network \mathcal{N} .

By the first theorem we prove the existence of the greatest regular fuzzy equivalence contained in a given n -tuple of fuzzy equivalences on the components of a multi-mode fuzzy network. We also show that it is the greatest solution, contained the same n -tuple of fuzzy equivalences, to a particular system of fuzzy relation equations.

Theorem 3.1 *Let $\mathcal{N} = (A_1, \dots, A_n, \mathcal{R})$ be a multi-mode fuzzy network, for each $l \in [1, n]$ let ξ_l denote an unknown taking values in $L^{A_l \times A_l}$, and let $(\alpha_1^0, \dots, \alpha_n^0)$ be a given n -tuple of fuzzy equivalences on A_1, \dots, A_n , respectively. Then the system of fuzzy relation equations*

$$\xi_j \circ R_i^{j,k} = R_i^{j,k} \circ \xi_k, \quad (j, k) \in J, i \in I_{j,k}, \quad (8)$$

$$\xi_j^{-1} \circ R_i^{j,k} = R_i^{j,k} \circ \xi_k^{-1}, \quad (j, k) \in J, i \in I_{j,k}, \quad (9)$$

has the greatest solution contained in $(\alpha_1^0, \dots, \alpha_n^0)$, which is the greatest regular fuzzy equivalence on \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

Proof. It is easy to see that the system (8)–(9) has at least one solution contained in $(\alpha_1^0, \dots, \alpha_n^0)$, the n -tuple $(\Delta_{A_1}, \dots, \Delta_{A_n})$ consisting of equality relations on A_1, \dots, A_n .

Let $\{(\alpha_1^t, \dots, \alpha_n^t)\}_{t \in T}$ be the family of all solutions to (8)–(9) contained in $(\alpha_1^0, \dots, \alpha_n^0)$, and for any $l \in [1, n]$ let

$$\alpha_l = \bigvee_{t \in T} \alpha_l^t.$$

Then for all $(j, k) \in J$ and $i \in I_{j,k}$ we have that,

$$\begin{aligned} \alpha_j \circ R_i^{j,k} &= \left(\bigvee_{t \in T} \alpha_j^t \right) \circ R_i^{j,k} = \bigvee_{t \in T} (\alpha_j^t \circ R_i^{j,k}) \\ &= \bigvee_{t \in T} (R_i^{j,k} \circ \alpha_k^t) = R_i^{j,k} \circ \left(\bigvee_{t \in T} \alpha_k^t \right) \\ &= R_i^{j,k} \circ \alpha_k, \end{aligned}$$

and similarly we prove that $\alpha_j^{-1} \circ R_i^{j,k} = R_i^{j,k} \circ \alpha_k^{-1}$, for all $(j, k) \in J$ and $i \in I_{j,k}$. Thus, $(\alpha_1, \dots, \alpha_n)$ is a solution to the system (8)–(9), and evidently, it is the greatest solution to this system contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

Further, it is easy to verify that $(\Delta_{A_1}, \dots, \Delta_{A_n})$, $(\alpha_1^{-1}, \dots, \alpha_n^{-1})$ and $(\alpha_1 \circ \alpha_1, \dots, \alpha_n \circ \alpha_n)$ are also solutions to (8)–(9). Since $(\alpha_1, \dots, \alpha_n)$ is the greatest solution to this system contained in $(\alpha_1^0, \dots, \alpha_n^0)$, we conclude that $\Delta_{A_l} \leq \alpha_l$, $\alpha_l^{-1} \leq \alpha_l$ and $\alpha_l \circ \alpha_l \leq \alpha_l$, which means that α_l is a fuzzy equivalence, for every $l \in [1, n]$, and therefore, the n -tuple $(\alpha_1, \dots, \alpha_n)$ is a regular fuzzy equivalence on the network \mathcal{N} .

Finally, every regular fuzzy equivalence on \mathcal{N} is a solution to (8)–(9), so we conclude that $(\alpha_1, \dots, \alpha_n)$ is the greatest regular fuzzy equivalence on \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$. \square

Note that regular fuzzy equivalences are n -tuples of fuzzy equivalences that are solutions to system (8). However, the greatest solution to this system, contained in a given n -tuple of fuzzy equivalences, is an n -tuple of fuzzy quasi-orders, but it is not necessarily an n -tuple of fuzzy equivalences. To ensure that this greatest solution is an n -tuple of fuzzy equivalences, the system (8) must be combined with the system (9), as was done in the previous theorem.

Next, for each $j \in [1, n]$ we set

$$\begin{aligned} \Lambda_j &= \{k \in [1, n] \mid (j, k) \in J\}, \\ P_j &= \{k \in [1, n] \mid (k, j) \in J\}. \end{aligned}$$

Moreover, for a given family \mathcal{F} of fuzzy relations, $\mathcal{L}(\mathcal{F})$ denotes the subalgebra of \mathcal{L} generated by all membership values taken by fuzzy relations from \mathcal{F} .

The following theorem gives a procedure for computing the greatest regular fuzzy equivalence contained in a given tuple of fuzzy equivalences.

Theorem 3.2 *Let $\mathcal{N} = (A_1, \dots, A_n, \mathcal{R})$ be a multi-mode fuzzy network, and let $(\alpha_1^0, \dots, \alpha_n^0)$ be a given n -tuple of fuzzy equivalences on A_1, \dots, A_n , respectively, and let $\{(\alpha_1^r, \dots, \alpha_n^r)\}_{r \in \mathbb{N}}$ be a descending sequence of n -tuples of fuzzy relations on A_1, \dots, A_n ,*

respectively, defined inductively as follows:

$$(\alpha_1^1, \dots, \alpha_n^1) = (\alpha_1^0, \dots, \alpha_n^0), \quad (10)$$

$$\begin{aligned} \alpha_j^{r+1} &= \alpha_j^r \wedge \bigwedge_{k \in \Lambda_j} \bigwedge_{i \in I_{j,k}} \left([(R_i^{j,k} \circ \alpha_k^r) / R_i^{j,k}] \wedge \right. \\ &\quad \left. \wedge [(R_i^{j,k} \circ (\alpha_k^r)^{-1}) / R_i^{j,k}]^{-1} \right) \wedge \\ &\quad \wedge \bigwedge_{k \in P_j} \bigwedge_{i \in I_{k,j}} \left([R_i^{k,j} \setminus (\alpha_k^r \circ R_i^{k,j})] \wedge \right. \\ &\quad \left. \wedge [R_i^{k,j} \setminus ((\alpha_k^r)^{-1} \circ R_i^{k,j})]^{-1} \right), \end{aligned} \quad (11)$$

for all $j \in [1, n]$, $r \in \mathbb{N}$. Then the following is true:

(a) *if there exists $s \in \mathbb{N}$ such that*

$$(\alpha_1^s, \dots, \alpha_n^s) = (\alpha_1^{s+1}, \dots, \alpha_n^{s+1}),$$

then the n -tuple $(\alpha_1^s, \dots, \alpha_n^s)$ is the greatest regular fuzzy equivalence on the fuzzy network \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$;

(b) *if A_1, \dots, A_n are finite sets and the subalgebra $\mathcal{L}(\mathcal{R} \cup \{\alpha_1^0, \dots, \alpha_n^0\})$ satisfies DCC, then the sequence $\{(\alpha_1^r, \dots, \alpha_n^r)\}_{r \in \mathbb{N}}$ is finite and there is $s \in \mathbb{N}$ so that $(\alpha_1^s, \dots, \alpha_n^s) = (\alpha_1^{s+1}, \dots, \alpha_n^{s+1})$.*

Proof. (a) Suppose that there exists $s \in \mathbb{N}$ such that $(\alpha_1^s, \dots, \alpha_n^s) = (\alpha_1^{s+1}, \dots, \alpha_n^{s+1})$. Then we have that

$$\alpha_j^s \leq (R_i^{j,k} \circ \alpha_k^s) / R_i^{j,k}$$

and

$$(\alpha_j^s)^{-1} \leq (R_i^{j,k} \circ (\alpha_k^s)^{-1}) / R_i^{j,k},$$

for all $j \in [1, n]$, $k \in \Lambda_j$ and $i \in I_{j,k}$, and also,

$$\alpha_j^s \leq R_i^{k,j} \setminus (\alpha_k^s \circ R_i^{k,j})$$

and

$$(\alpha_j^s)^{-1} \leq R_i^{k,j} \setminus ((\alpha_k^s)^{-1} \circ R_i^{k,j}),$$

for all $j \in [1, n]$, $k \in P_j$ and $i \in I_{k,j}$. According to the residuation property, this is equivalent to

$$\alpha_j^s \circ R_i^{j,k} \leq R_i^{j,k} \circ \alpha_k^s$$

and

$$(\alpha_j^s)^{-1} \circ R_i^{j,k} \leq R_i^{j,k} \circ (\alpha_k^s)^{-1},$$

for all $j \in [1, n]$, $k \in \Lambda_j$ and $i \in I_{j,k}$, and also,

$$R_i^{k,j} \circ \alpha_j^s \leq \alpha_k^s \circ R_i^{k,j}$$

and

$$R_i^{k,j} \circ (\alpha_j^s)^{-1} \leq (\alpha_k^s)^{-1} \circ R_i^{k,j},$$

for all $j \in [1, n]$, $k \in P_j$ and $i \in I_{k,j}$. It is not hard to verify that the system consisting of the previous four types of fuzzy relation inequalities is equivalent to the system (8)–(9), and consequently, the n -tuple $(\alpha_1^s, \dots, \alpha_n^s)$ is a solution to the system (8)–(9). Evidently, it is contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

Let $(\alpha_1, \dots, \alpha_n)$ be an arbitrary solution to the system (8)–(9) contained in $(\alpha_1^0, \dots, \alpha_n^0)$. Suppose that $(\alpha_1, \dots, \alpha_n) \leq (\alpha_1^r, \dots, \alpha_n^r)$, for some $r \in \mathbb{N}^0$.

Then for each $j \in [1, n]$, $k \in \Lambda_j$ and $i \in I_{j,k}$ we have that

$$\alpha_j \circ R_i^{j,k} \leq R_i^{j,k} \circ \alpha_k \leq R_i^{j,k} \circ \alpha_k^r,$$

whence

$$\alpha_j \leq (R_i^{j,k} \circ \alpha_k^r) / R_i^{j,k},$$

and similarly,

$$(\alpha_j)^{-1} \leq (R_i^{j,k} \circ (\alpha_k^r)^{-1}) / R_i^{j,k},$$

which implies that

$$\begin{aligned} \alpha_j \leq \alpha_j^r \wedge \bigwedge_{k \in \Lambda_j} \bigwedge_{i \in I_{j,k}} & \left([(R_i^{j,k} \circ \alpha_k^r) / R_i^{j,k}] \wedge \right. \\ & \left. \wedge [(R_i^{j,k} \circ (\alpha_k^r)^{-1}) / R_i^{j,k}]^{-1} \right). \end{aligned}$$

In the same way we show that

$$\begin{aligned} \alpha_j \leq \alpha_j^r \wedge \bigwedge_{k \in P_j} \bigwedge_{i \in I_{k,j}} & \left([R_i^{k,j} \setminus (\alpha_k^r \circ R_i^{k,j})] \wedge \right. \\ & \left. \wedge [R_i^{k,j} \setminus ((\alpha_k^r)^{-1} \circ R_i^{k,j})]^{-1} \right). \end{aligned}$$

Now, according to (11) we have that $\alpha_j \leq \alpha_j^{r+1}$, for each $j \in [1, n]$, and by induction we conclude that $\alpha_j \leq \alpha_j^r$, for all $j \in [1, n]$ and $r \in \mathbb{N}^0$.

Consequently, $\alpha_j \leq \alpha_j^s$, for each $j \in [1, n]$, which means that the n -tuple $(\alpha_1^s, \dots, \alpha_n^s)$ is the greatest solution to the system (8)–(9) contained in the n -tuple $(\alpha_1^0, \dots, \alpha_n^0)$. Now, according to Theorem 3.1, we conclude that $(\alpha_1^s, \dots, \alpha_n^s)$ is the greatest regular fuzzy equivalence on \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

(b) Suppose that A_1, \dots, A_n are finite sets and the subalgebra $\mathcal{L}(\mathcal{R} \cup \{\alpha_1^0, \dots, \alpha_n^0\})$ satisfies DCC.

For each $l \in [1, n]$ and all pairs $(a_l, a'_l) \in A_l \times A_l$, we have that $\{\alpha_l^r(a_l, a'_l)\}_{r \in \mathbb{N}}$ is a descending sequence in $\mathcal{L}(\mathcal{R} \cup \{\alpha_1^0, \dots, \alpha_n^0\})$. By the hypothesis, this sequence stabilizes, and since there are finitely many such sequences, we conclude that there exists $s \in \mathbb{N}$ such that all these sequences stabilize after s steps. Therefore, the sequence $\{(\alpha_1^r, \dots, \alpha_n^r)\}_{r \in \mathbb{N}}$ is finite and $(\alpha_1^s, \dots, \alpha_n^s) = (\alpha_1^{s+1}, \dots, \alpha_n^{s+1})$ for $s \in \mathbb{N}$ whose existence has been established above.

This completes the proof of the theorem. \square

By Theorem 3.2, the greatest regular fuzzy equivalence contained in the given n -tuple $(\alpha_1^0, \dots, \alpha_n^0)$ of fuzzy equivalences can be computed as follows. We start from this n -tuple and build a descending sequence of n -tuples of fuzzy equivalences by means of the formula (11), and simultaneously we check whether two subsequent members of the sequence are equal. The procedure terminates as soon as we find the first pair of equal consecutive members of the sequence, and in this case, the last computed n -tuple is the greatest regular fuzzy equivalence contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

However, in general, the above described procedure do not necessarily terminate in a finite number of steps, and Theorem 3.2 provides a sufficient condition under which it will terminate, when the subalgebra $\mathcal{L}(\mathcal{R} \cup \{\alpha_1^0, \dots, \alpha_n^0\})$ satisfies the descending

chain condition. In particular, this condition is fulfilled if \mathcal{L} is a locally finite algebra, that is, if every finitely generated subalgebra of \mathcal{L} is finite. The most widely used locally finite structures are the Boolean structure and the Gödel structure. For more information on local finiteness in t-norm based structures we refer to the recent paper [13].

In cases when the above procedure does not terminate in a finite number of steps and can not be used to efficiently compute the greatest regular fuzzy equivalence on the network \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$, it is possible to modify this procedure to compute the greatest regular crisp equivalence contained in $(\alpha_1^0, \dots, \alpha_n^0)$. This modified procedure is provided by the following theorem.

Theorem 3.3 *Let $\mathcal{N} = (A_1, \dots, A_n, \mathcal{R})$ be a multi-mode fuzzy network, and let $(\alpha_1^0, \dots, \alpha_n^0)$ be a given n -tuple of fuzzy equivalences on A_1, \dots, A_n , respectively, and let $\{(\varrho_1^r, \dots, \varrho_n^r)\}_{r \in \mathbb{N}}$ be a descending sequence of n -tuples of crisp relations on A_1, \dots, A_n , respectively, defined inductively as follows:*

$$(\varrho_1^1, \dots, \varrho_n^1) = ([\alpha_1^0]^c, \dots, [\alpha_n^0]^c), \quad (12)$$

$$\varrho_j^{r+1} = \varrho_j^r \wedge [\Phi(\varrho_j^r)]^c, \quad (13)$$

where

$$\begin{aligned} \Phi(\varrho_j^r) = \bigwedge_{k \in \Lambda_j} \bigwedge_{i \in I_{j,k}} & \left([(R_i^{j,k} \circ \alpha_k^r) / R_i^{j,k}] \wedge \right. \\ & \left. \wedge [(R_i^{j,k} \circ (\alpha_k^r)^{-1}) / R_i^{j,k}]^{-1} \right) \wedge \\ & \bigwedge_{k \in P_j} \bigwedge_{i \in I_{k,j}} \left([R_i^{k,j} \setminus (\alpha_k^r \circ R_i^{k,j})] \wedge \right. \\ & \left. \wedge [R_i^{k,j} \setminus ((\alpha_k^r)^{-1} \circ R_i^{k,j})]^{-1} \right), \end{aligned}$$

for all $j \in [1, n]$, $r \in \mathbb{N}$.

Then there exists $s \in \mathbb{N}$ such that

$$(\varrho_1^s, \dots, \varrho_n^s) = (\varrho_1^{s+1}, \dots, \varrho_n^{s+1}),$$

and the n -tuple $(\varrho_1^s, \dots, \varrho_n^s)$ is the greatest regular crisp equivalence on the fuzzy network \mathcal{N} contained in $(\alpha_1^0, \dots, \alpha_n^0)$.

The proof of this theorem is similar to the proof of Theorem 3.2 and it will be omitted. The reader can also take a look at Proposition 5.8 [15].

4. Concluding remarks

In this article we dealt with regular fuzzy equivalences on multi-mode fuzzy networks because such equivalences are generally recognized in the social network analysis as a powerful mean for identifying, describing and understanding the social positions. The subject of our further research will be certain more general types of fuzzy equivalences and fuzzy quasi-orders, which will also be obtained as solutions to particular systems of fuzzy relation equations and inequalities, and which we will also try to

apply in the study of the social positions. As noted in [5], which equivalence relation is interesting to consider depends on the problem at hand, and it is likely necessary to consider several different equivalence relations for a given network, in order to understand it completely. For example, in the mentioned paper [5] the authors applied the so-called simulation equivalences in the analysis of a network that represents the communication between a group of terrorists with different social positions, and they have shown that this kind of equivalences may help to identify some social positions that cannot be recognized using regular equivalences.

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